

# CS 153 Lab6

Kishore Kumar Pusukuri



# Outline

Mutex vs Condition Variables  
Semaphores

- ▶ Unlocking and locking mutex leads spinning or polling, wastes CPU time.
- ▶ We could sleep for some amount of time, but we do not know how long to sleep.
- ▶ A mutex is for locking and a condition variable is for waiting.
- ▶ Study the program prodcons\_1.c

```
#include <pthread.h>
```

```
int pthread_cond_wait(pthread_cond_t *cptr,  
pthread_mutex_t *mptr);  
int pthread_cond_signal(pthread_cond_t *cptr);
```

Both return: 0 if OK, positive Exxx value on error.

A conditional variable is a variable of type  
pthread\\_cond\\_t.

# Conditional variable functions

- ▶ A mutex is always associated with a condition variable.
- ▶ When we call `pthread_cond_wait` to wait for some condition to be true, we specify the address of the condition variable and the address of the associated mutex.
- ▶ The term “signal” in the second function’s name does NOT refer to a `SIGxxx` signal.
- ▶ Normally, `pthread_cond_signal` awakes one thread that is waiting on the condition variable.
- ▶ Study the program `prodcons_2.c`

# Review

- ▶ Mutexes are used to protect critical regions of code, so that only one thread at a time is executing within the critical region.
- ▶ Sometimes a thread obtains a mutex lock and then discovers that it needs to wait for some condition to be true. When this happens, the thread waits on a condition variable.
- ▶ A condition variable is always associated with a mutex.
- ▶ The `pthread_cond_wait` function that puts the thread to sleep unlocks the mutex before putting the thread to sleep and relocks the mutex before waking up the thread at some later time.
- ▶ The condition variable is signaled by some other thread, and that signaling thread has the option of waking up one thread using `pthread_cond_signal` or all the threads that are waiting for the condition to be true (using `pthread_cond_broadcast`). 



# Why semaphores ?

- ▶ Mutexes and conditional variables can always be used to synchronize the various threads within a process.
- ▶ Posix also allows a mutex or conditional variable to be used for synchronization between multiple processes, “if the mutex and conditional variable is stored in memory that is shared between processes” .
- ▶ A semaphore is a primitive used to provide synchronization between various processes or between the various threads in a given process.

Types :

- ▶ Posix named semaphores.
- ▶ Posix memory-based semaphores.
- ▶ System V Semaphores.

# Posix named semaphores

Consider binary semaphores :

- ▶ Posix semaphores are identified by names that might correspond to pathnames in the file system.
- ▶ Three operations that a process can perform on a semaphore :
  - ▶ Create a semaphore.
  - ▶ Wait for a semaphore.
  - ▶ Post a semaphore.

wait operation:

```
while (semaphore_value <= 0)
    ; /* wait; i.e., block the thread or process */
    semaphore_value--; /* we have the semaphore */
```

post/signal operation:

```
semaphore_value++;
```

Comparision of mutex and binary semaphore  
to solve mutual exclusion problem:

initialize mutex;	initialize semaphore sem to 1;
pthread_mutex_lock(&mutex)	sem_wait(&sem)
critical region	critical region
pthread_mutex_unlock(&mutex)	sem_post(&sem);

Semaphore functions:

```
#include <semaphore.h>
```

```
sem_t *sem_open(const char *name, int oflag,...  
                /*mode_t mode, unsigned int value */ );
```

Somewhat similar to syntax of file functions.

```
int sem_close(sem_t *sem);
```

```
int sem_wait (sem_t *sem);
```

```
int sem_post(sem_t *sem);
```

..and several other operations....

The above return: 0 if OK, -1 on error.

How to attack Project 2?



Thaaaank  
Yooooouuu!!!