



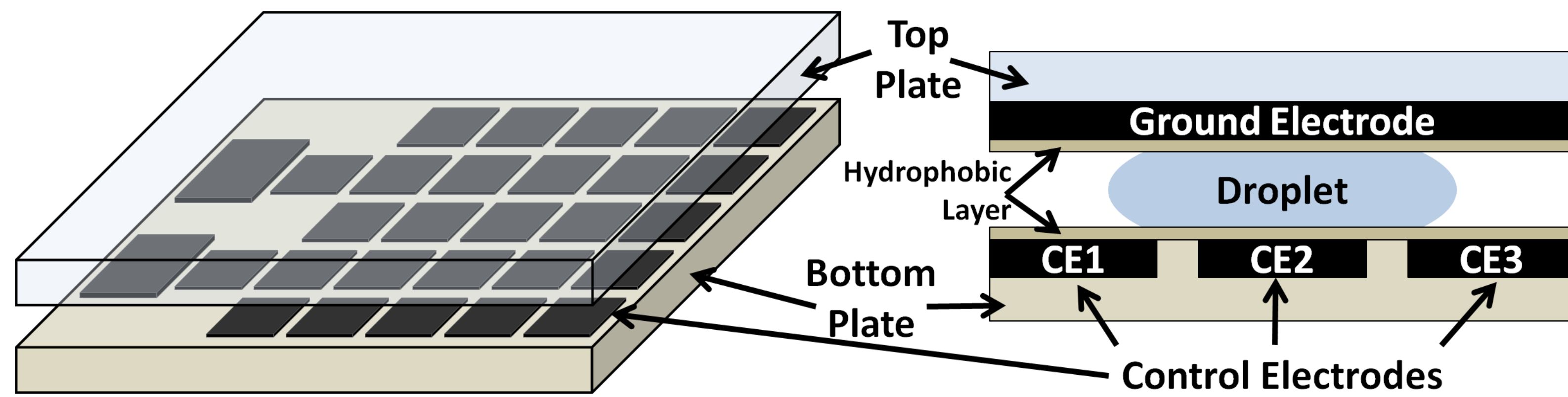
Path Scheduling on Digital Microfluidic Biochips

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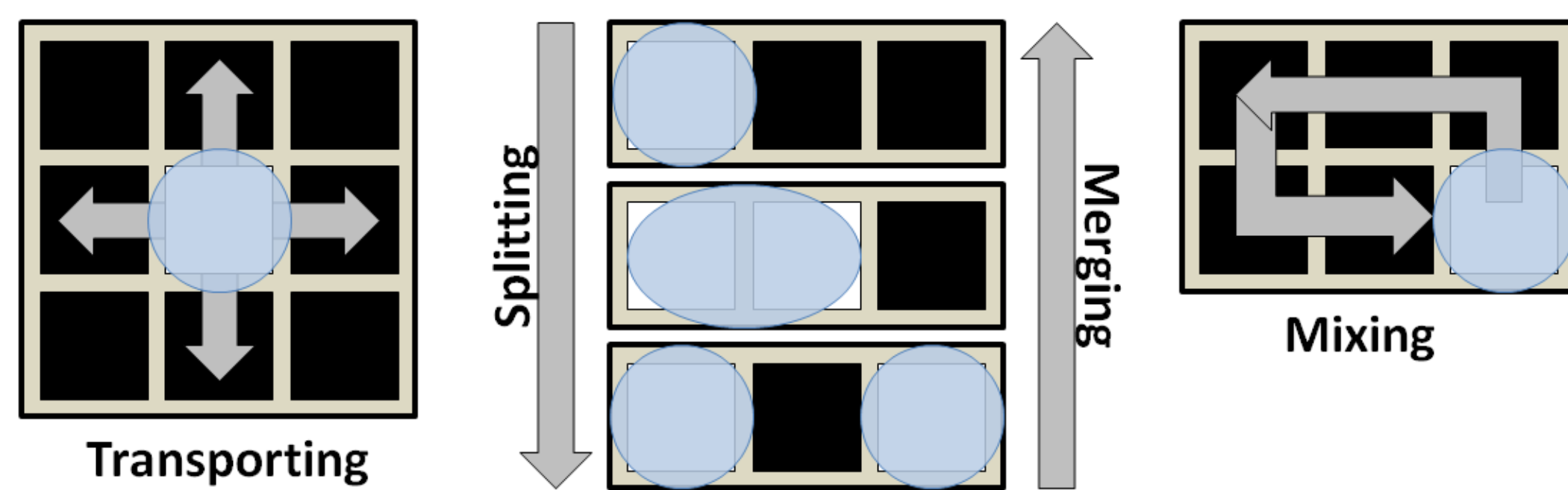


Digital Microfluidic Technology

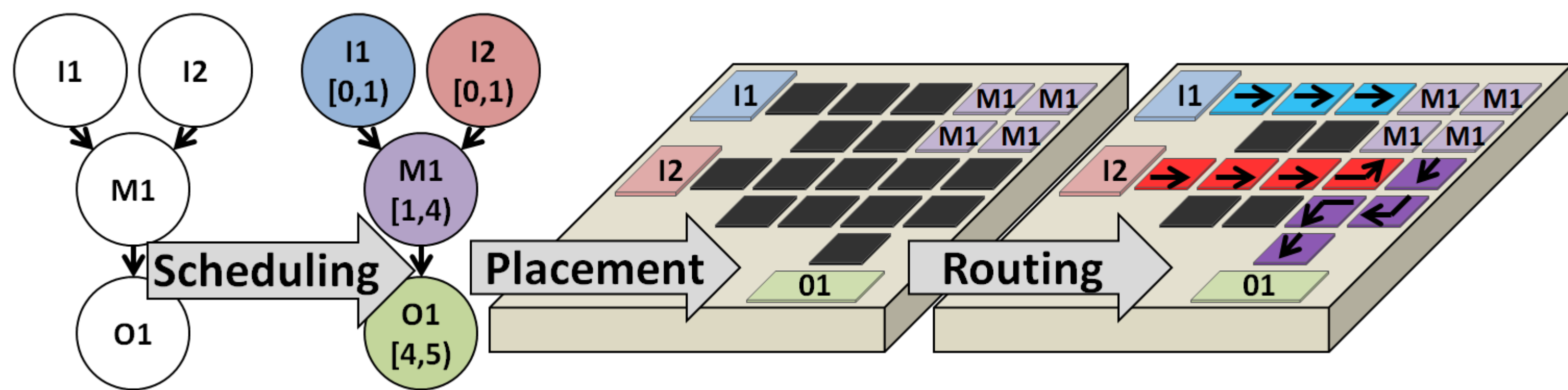
Digital Microfluidic Biochips (DMFBs) are an emerging “lab-on-a-chip (LoC)” technology that perform biochemical reactions by operating on fluidic droplets on the scale of nano-liters.



Microfluidic Array



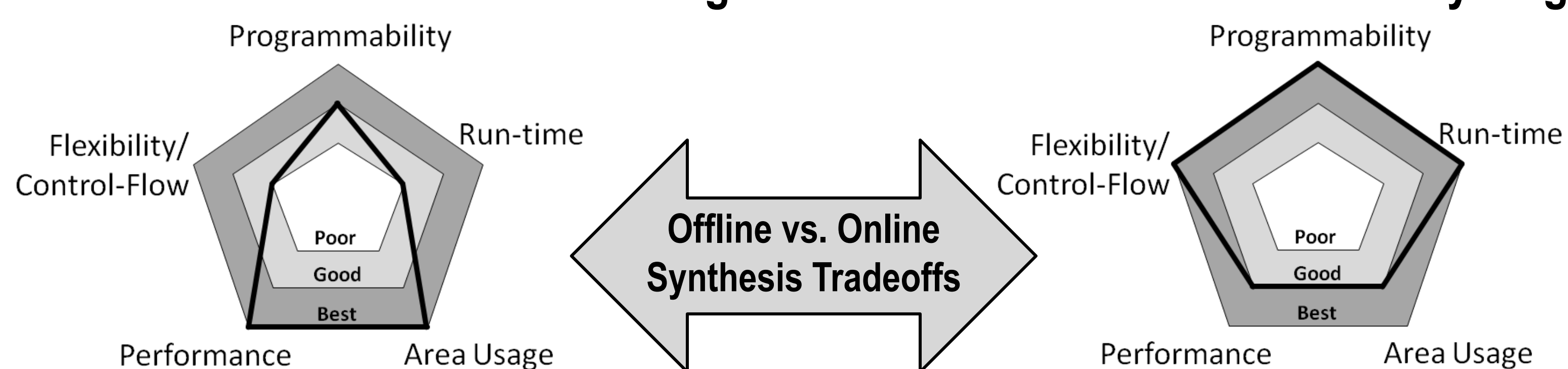
Basic Microfluidic Operations



Microfluidic Synthesis Flow

High-Level Motivation and Problem

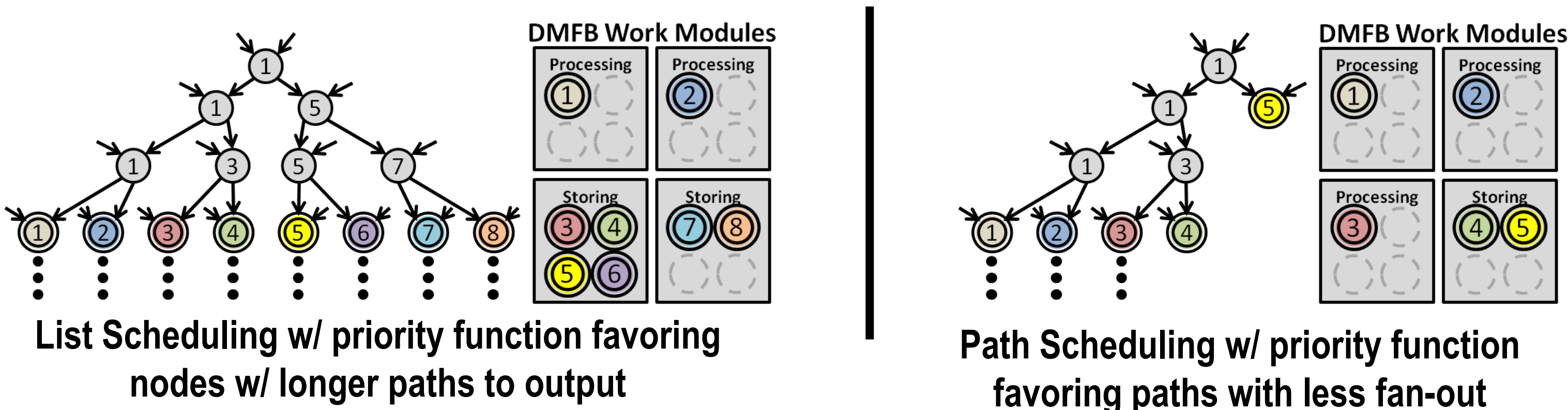
- Goal:** Online synthesis
- Motivation:** To enable DMFB programmability and new features in the areas of control-flow and live-feedback
- Example:** A control-flow graph which can dynamically decide which assay to run next based on live feedback from the DMFB:
- Problem:** Past offline synthesis methods are computationally complex, which would add significant amounts of time to the assay length



- Offline Assay Compilation**
- Solution:** Develop synthesis methods that yield good results in little time so the overall assay length is kept short
- Online Assay Interpretation**

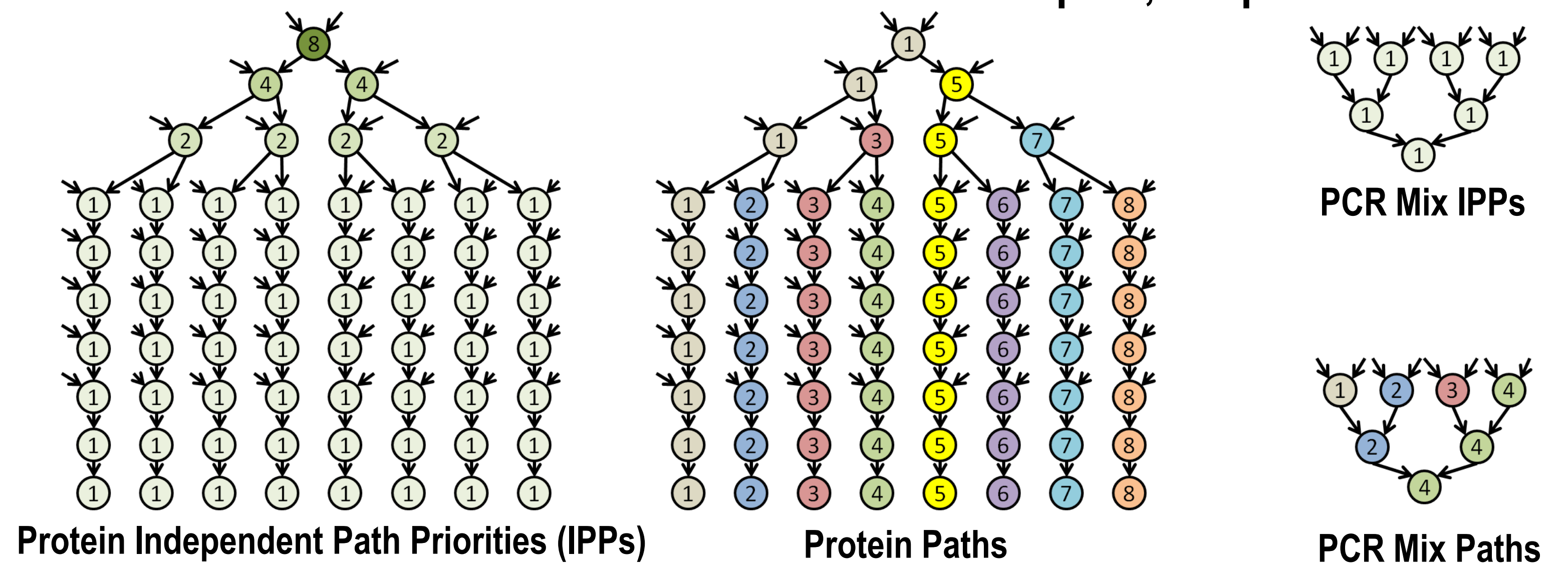
Minimizing Storage

- Traditional computing has “infinite” storage, as far as program is concerned
- DMFBs have limited storage since the same cells used for operations must also be used for storage if droplets are not ready to be operated on
- The order in which we schedule nodes can affect the amount of useful work the DMFB is performing:



Path Scheduling Algorithm

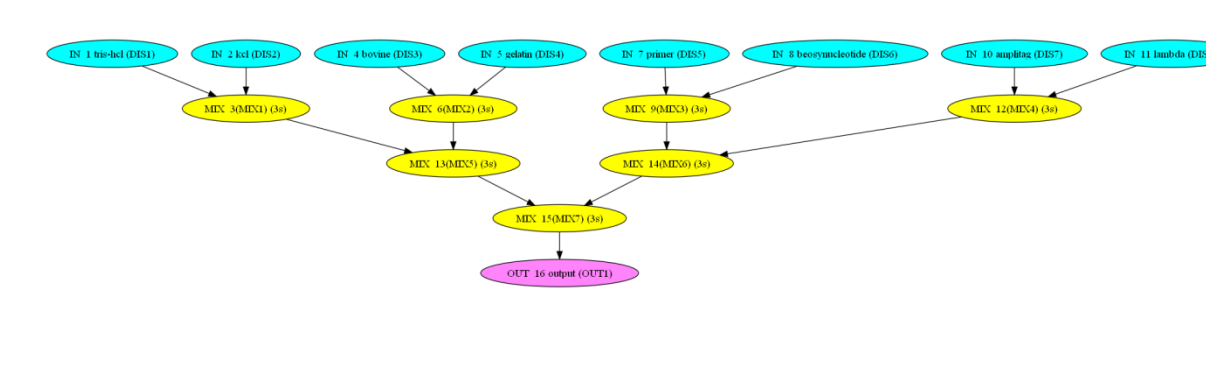
- Compute schedule for an entire path at a time (instead of node at a time)
- Start with a path-leader
 - Initially a node with only dispense parents
 - If there is a branch on the path:
 - Reserve resources for the node with the lowest IPP; continue down path
 - All other branch nodes are added to the list of path leaders for later
- Once resources have been reserved for an entire path, the path is scheduled



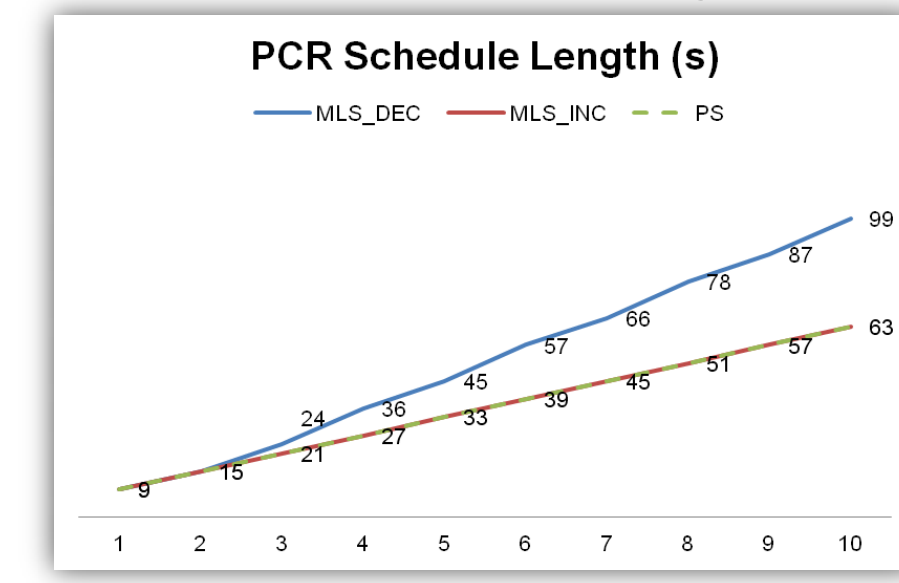
Results

PCR Mixing Stage Assay

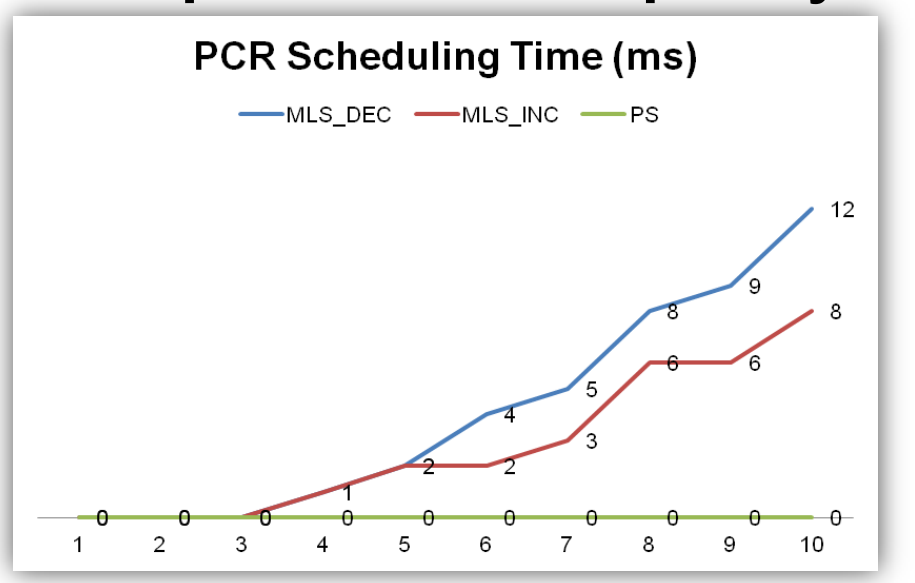
- Equivalent results in slightly less time



Solution Quality

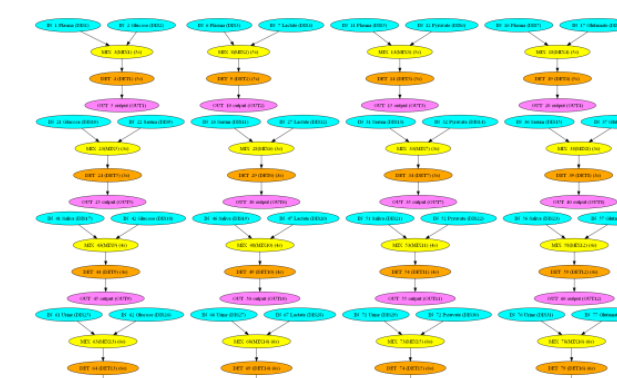


Computational Complexity

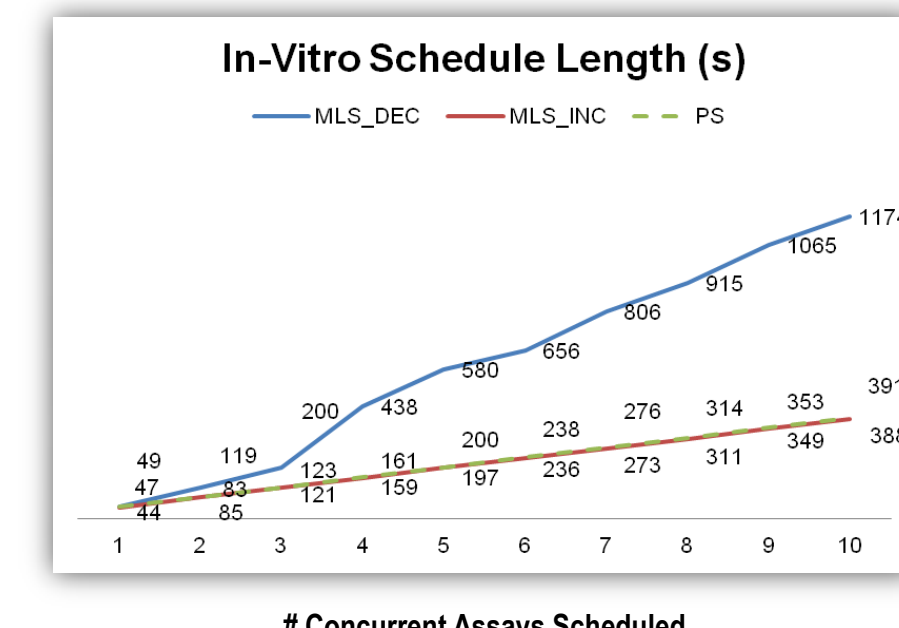


In-vitro Assay

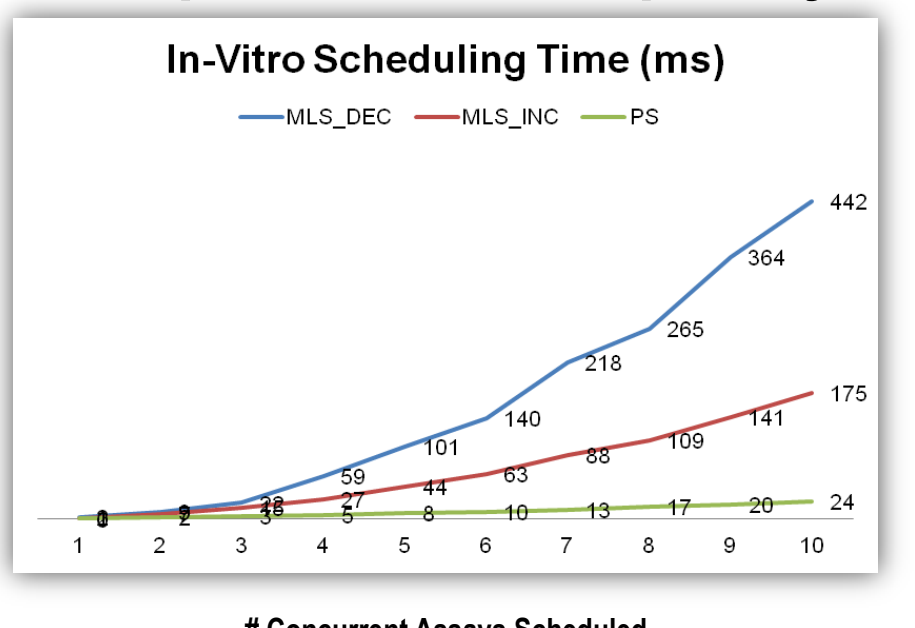
- PS 2-3 time-steps slower than MLS_INC
- Input inefficiency (details in paper)
- PS 3x to 7x faster



Solution Quality

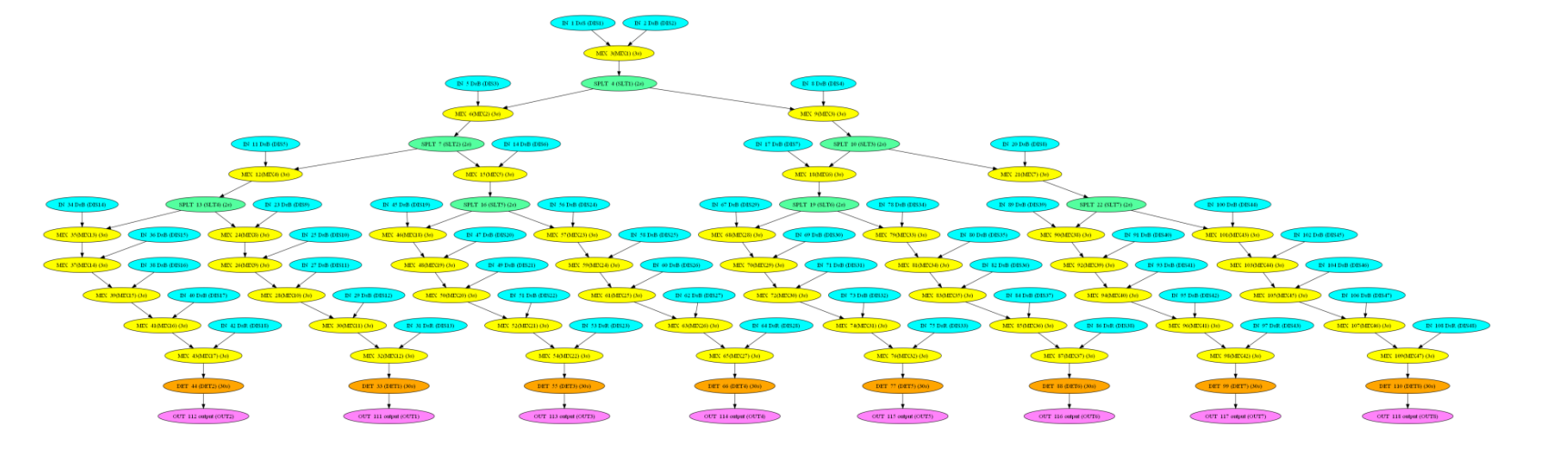


Computational Complexity

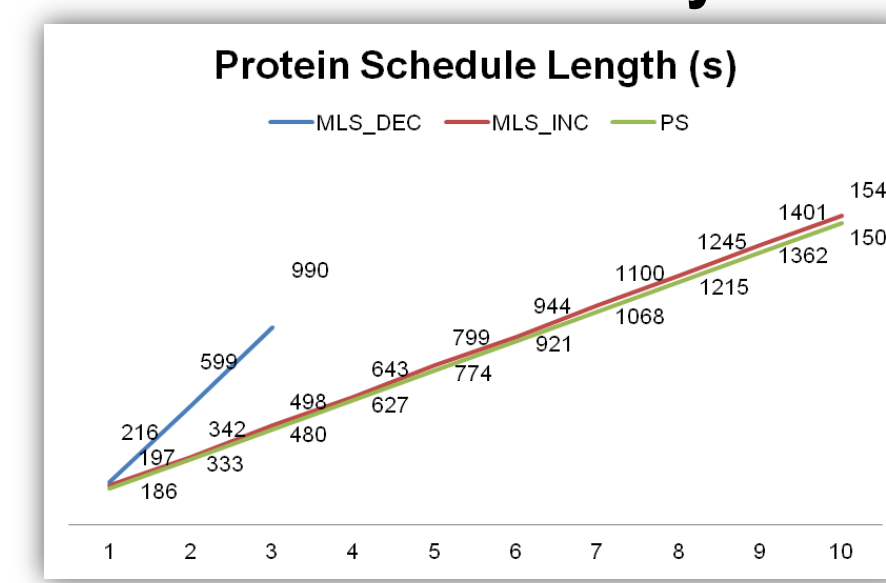


Protein Assay

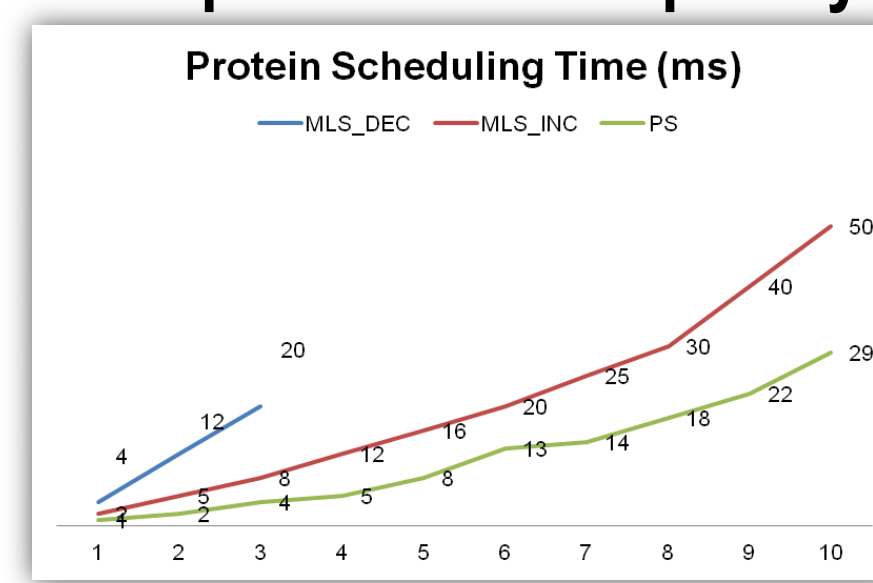
- MLS_DEC cannot schedule many DAGs b/c of storage inefficiency
- PS chops dozens of time-steps off schedule
- PS runs in about 1/3 the time of MLS_INC
- PS uses the fewest modules for storage
- As number of modules increases, solution quality converges



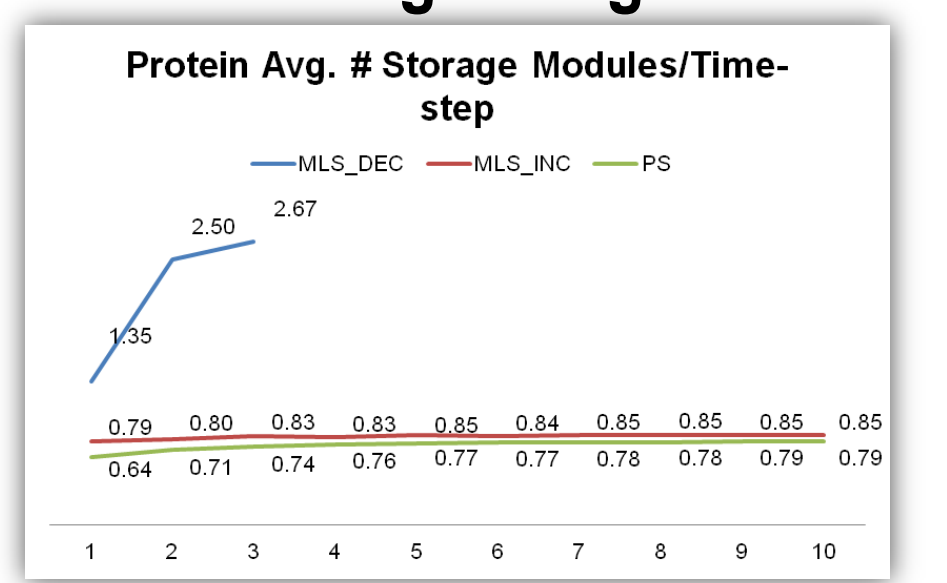
Solution Quality



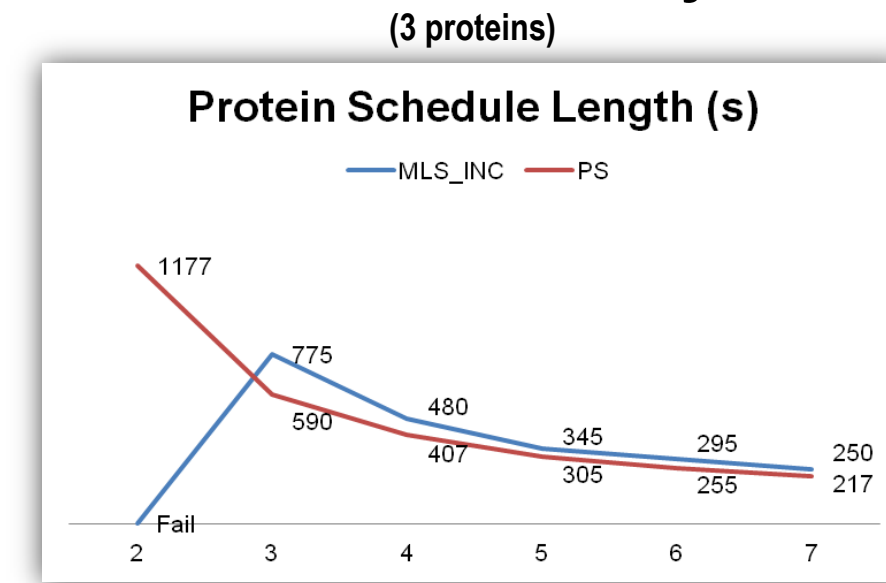
Computational Complexity



Storage Usage

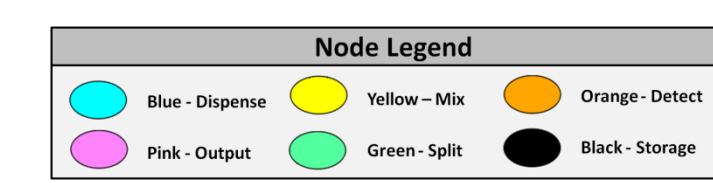


Solution Quality



Scheduled/Placed Protein:

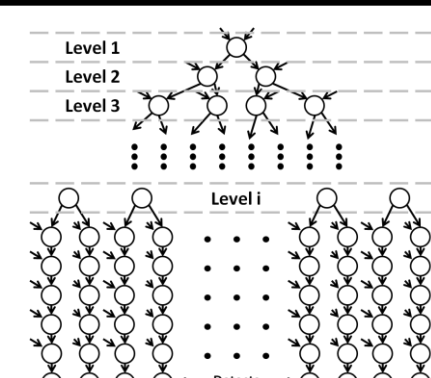
(Black node = droplet being stored for a # of time-steps in 1 module)



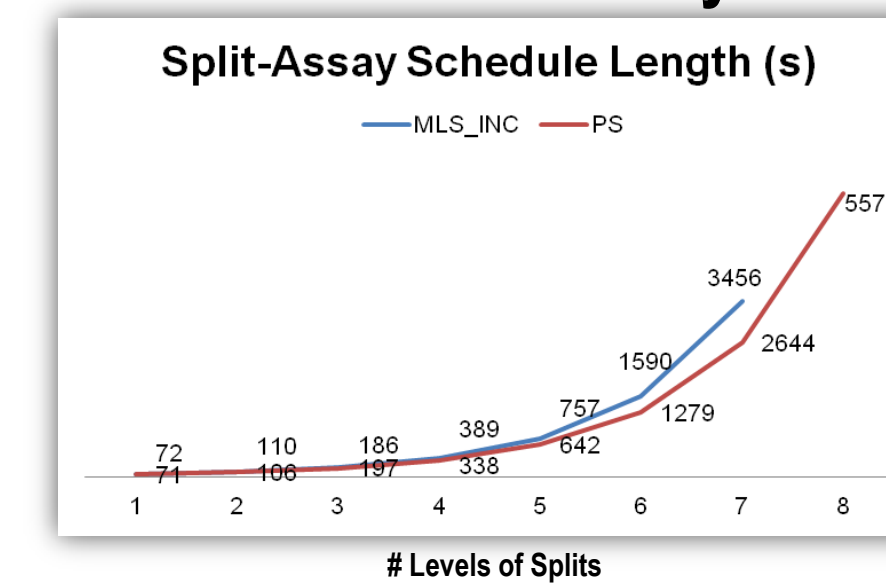
MLS_DEC MLS_INC PS

Split-Level Protein Assay

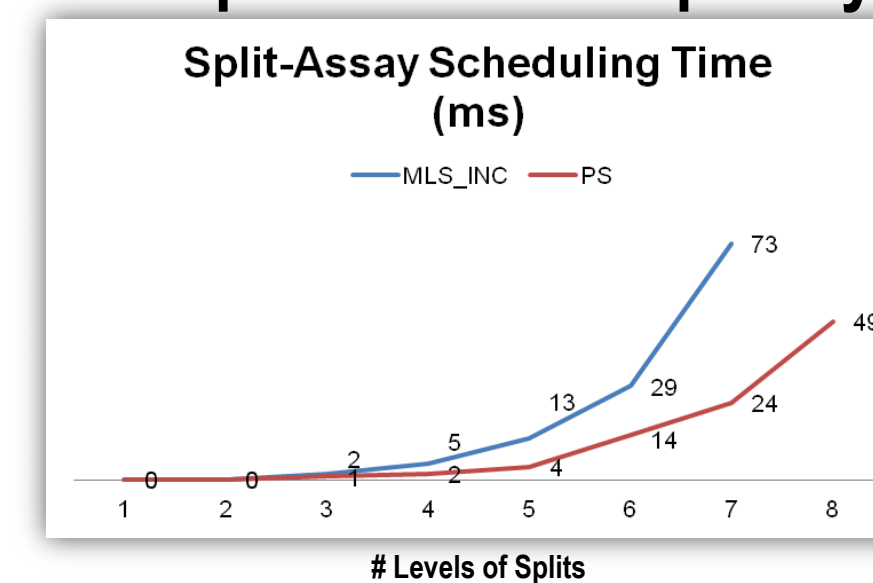
- PS saves hundreds of seconds as fan-out increases
- PS saves several dozen milliseconds of computation time
- MLS Fails at 8 splits



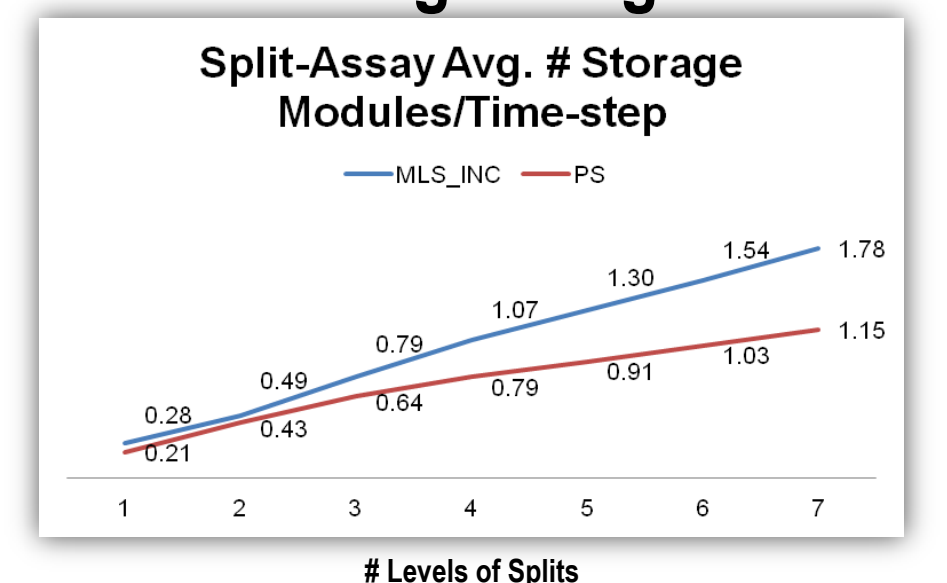
Solution Quality



Computational Complexity



Storage Usage



Contact

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